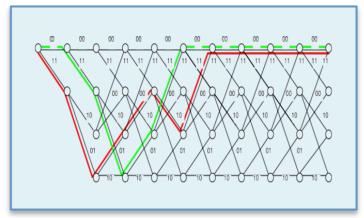
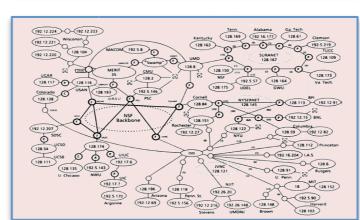


rf (freq. domain)





# INTRODUCTION TO EECS II DIGITAL COMMUNICATION

SYSTEMS

#### 6.02 Fall 2014 Lecture #23

- Reliable transport
  - Sliding-window protocol
  - Analysis of sliding-window

#### Unanswered questions

(about packet-switched networks)

#### How do nodes determine routes to every other node?

Nodes determine routes via either link-state, distance-vector, or path-vector routing

#### How do nodes route around link failures?

Routing protocols will **eventually converge**, but experience different problems along the way (routing loops, counting-to-infinity, etc.)

## How do nodes communicate reliably given that the network is best-effort?

Nodes can use a stop-and-wait protocol

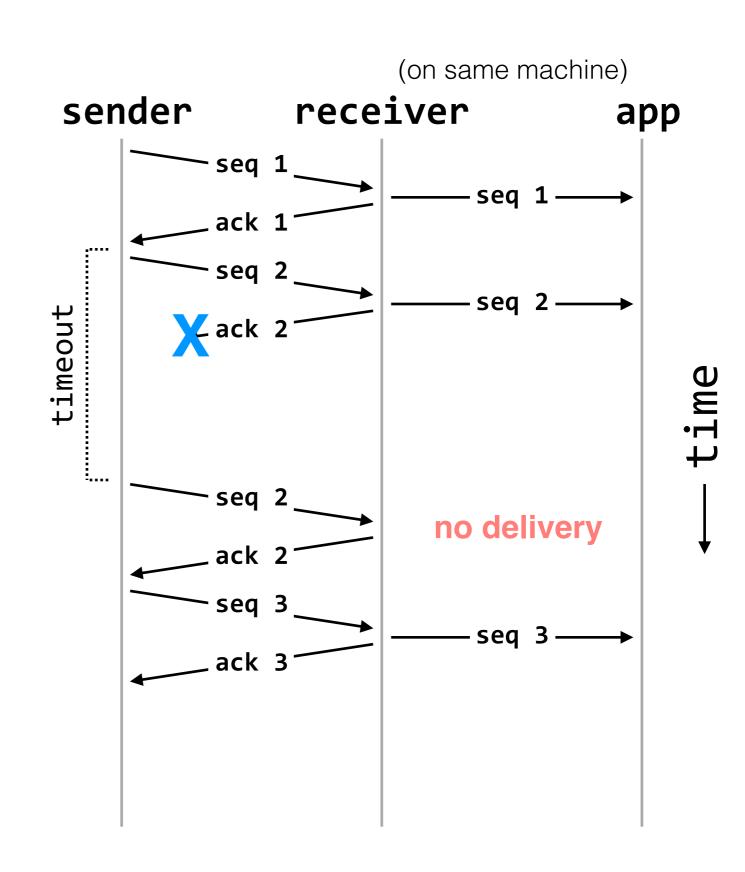
#### Recap: Stop-and-wait Protocol

#### At sender:

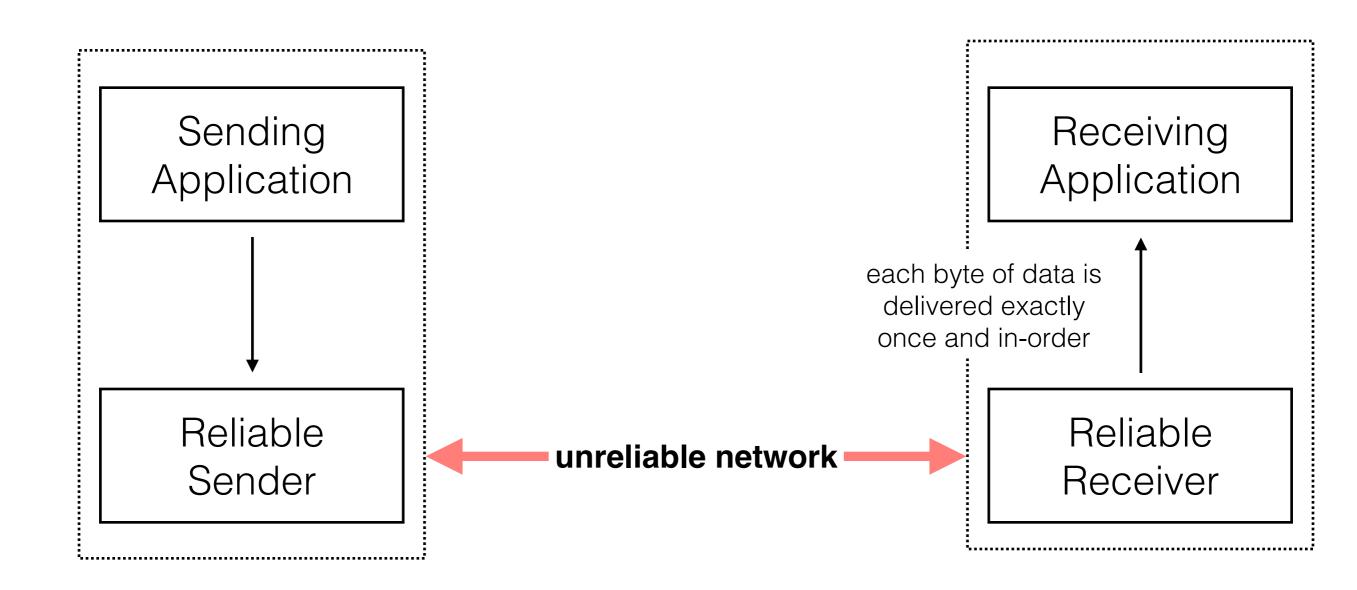
- Send a packet, keep track of its sequence number
- When an ACK is received for that packet, increment the stored sequence number and repeat
- If an ACK for the outstanding packet hasn't been received after timeout seconds, retransmit the packet

#### At receiver:

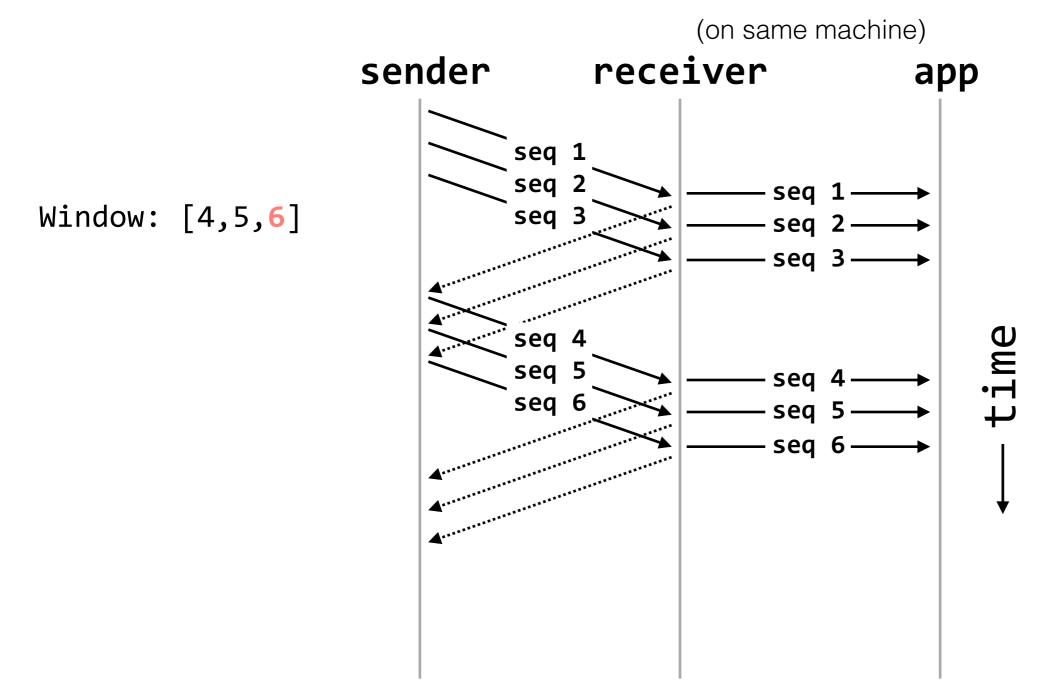
- Upon receipt of packet k, send an ACK for k
- If k is greater than the last received sequence number, deliver packet to app



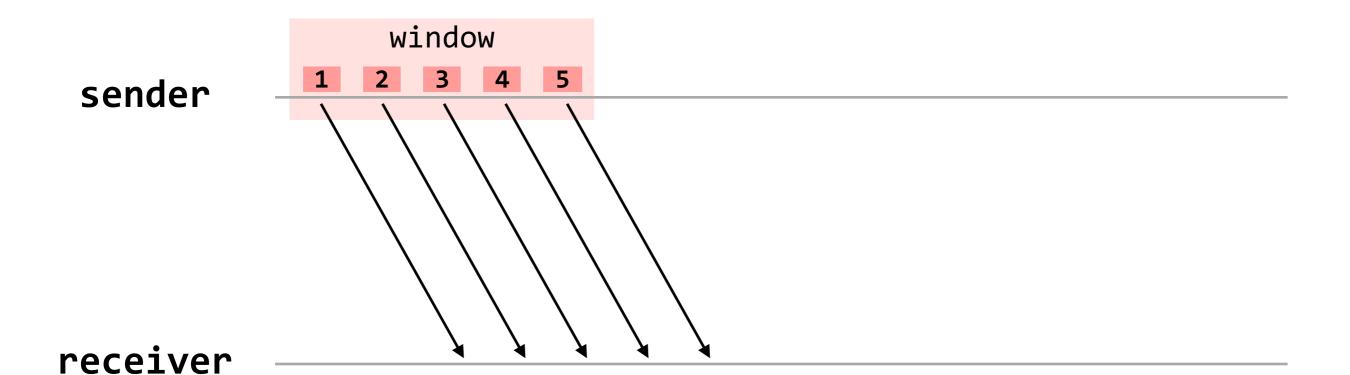
#### Reliable Communication

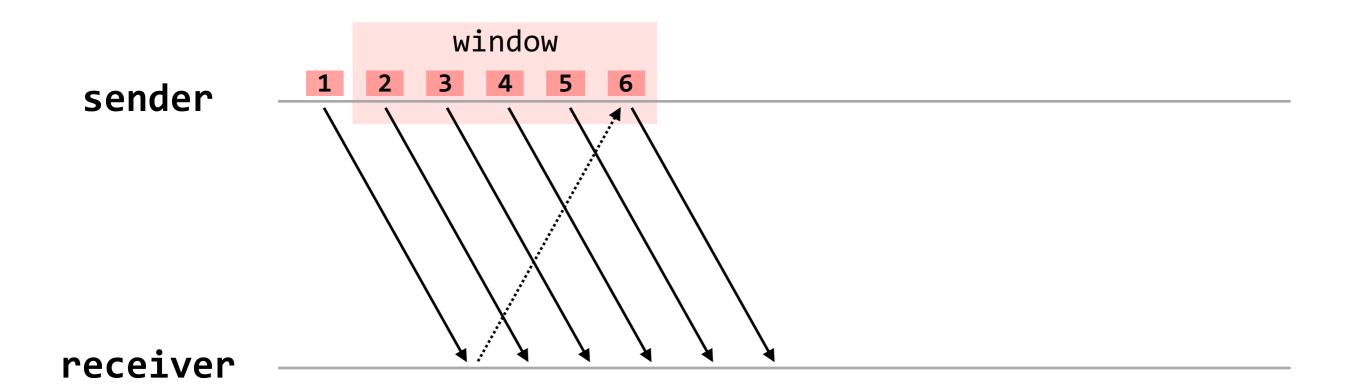


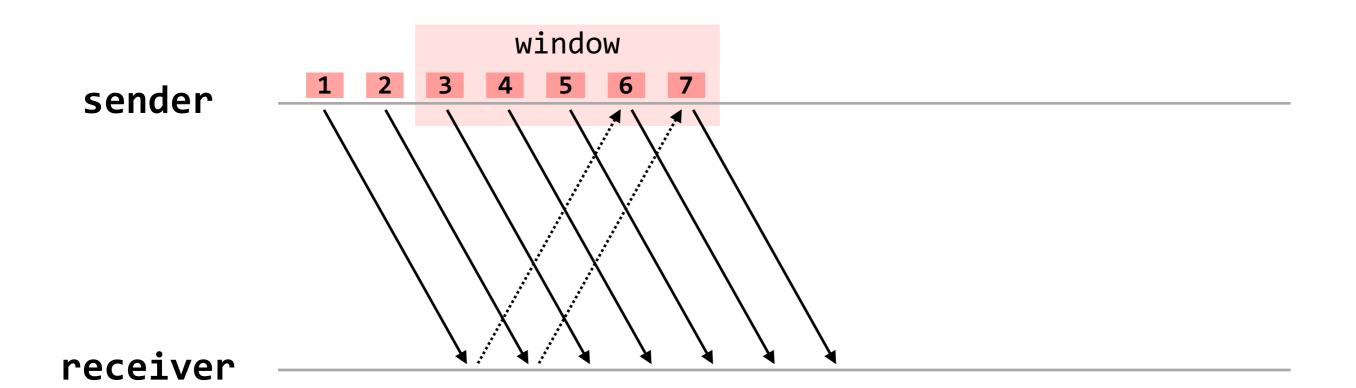
today's goal: develop a reliable transport protocol that gets better utilization than the stop-and-wait protocol

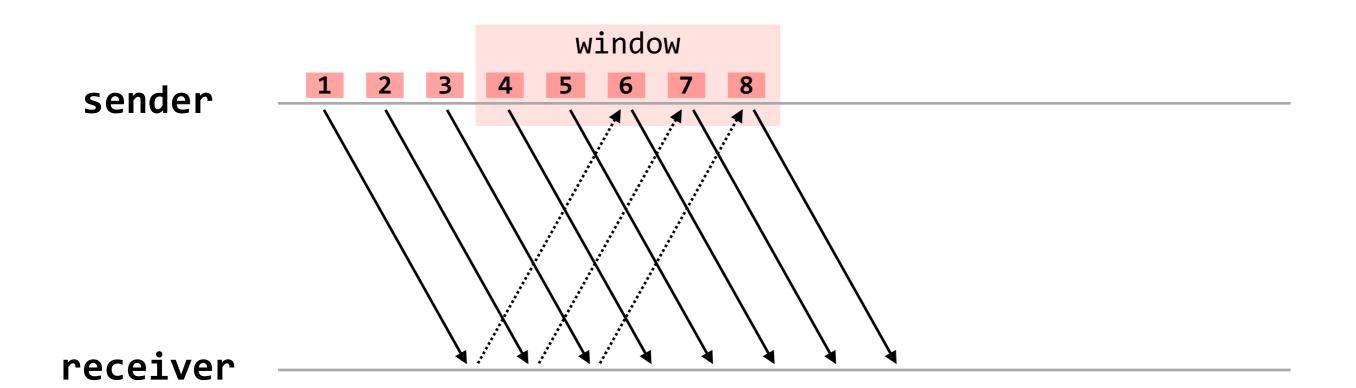


basic idea: send a new packet whenever an ACK is received, allowing no more than W outstanding packets at a time









	(On Same n	iaci ii ie)
sender	receiver	арр

window: [1,2,3,4,5]

seq 1

seq 2

seq 3

seq 4

seq 5

(on same machine)

sender receiver app

Window: [1,2,3,4,5] seq 1 seq seq

(on same machine)

sender receiver app Window: [2,3,4,5,6] seq 1 seq seq seq 6

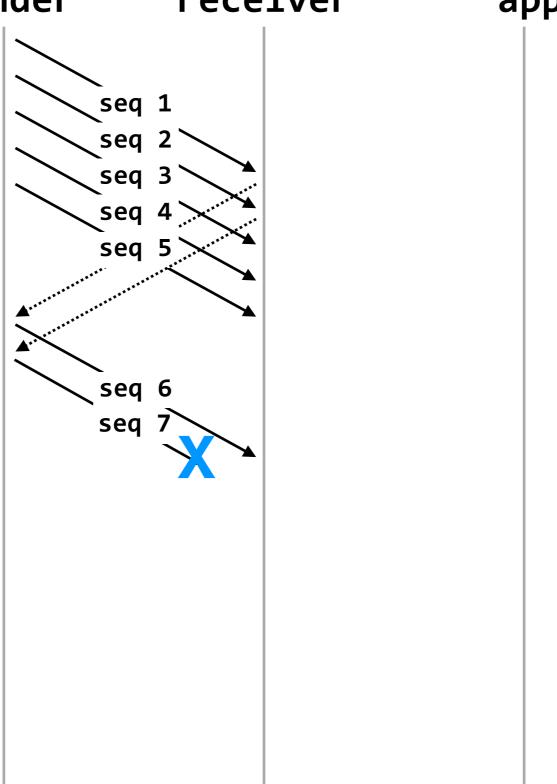
(on same machine)

sender receiver app Window: [2,3,4,5,6] seq 1 seq seq 6

(on same machine)

sender receiver app

Window: [3,4,5,6,7]

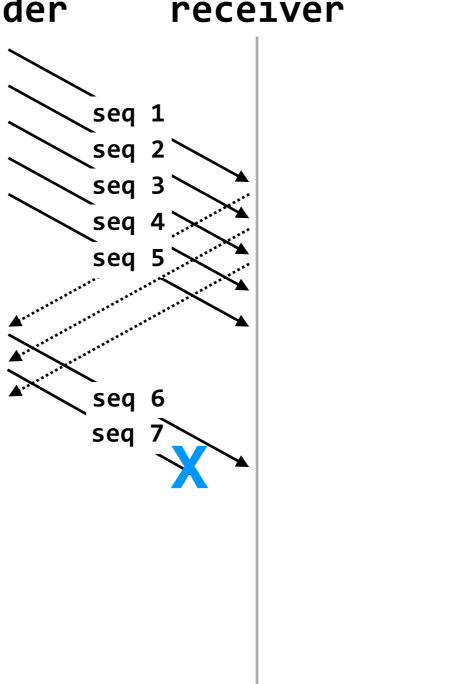


(on same machine)

app

sender receiver

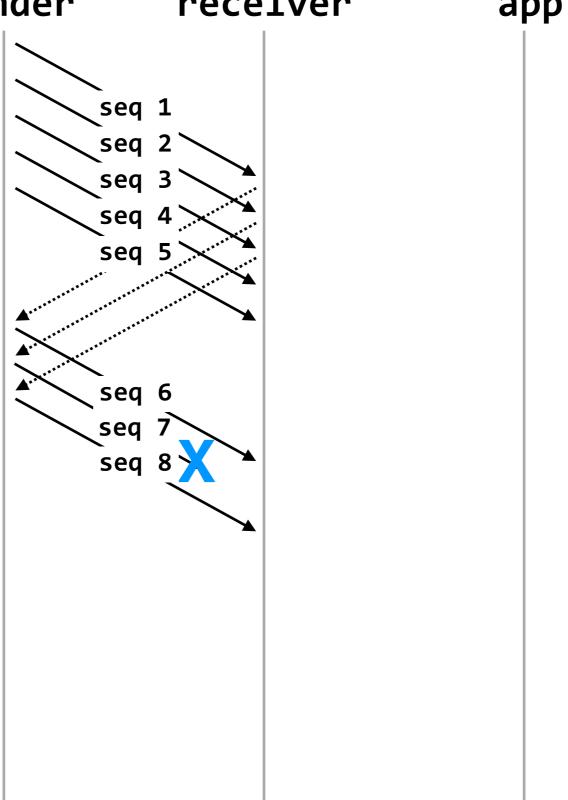
Window: [3,4,5,6,7]



(on same machine)

sender receiver app

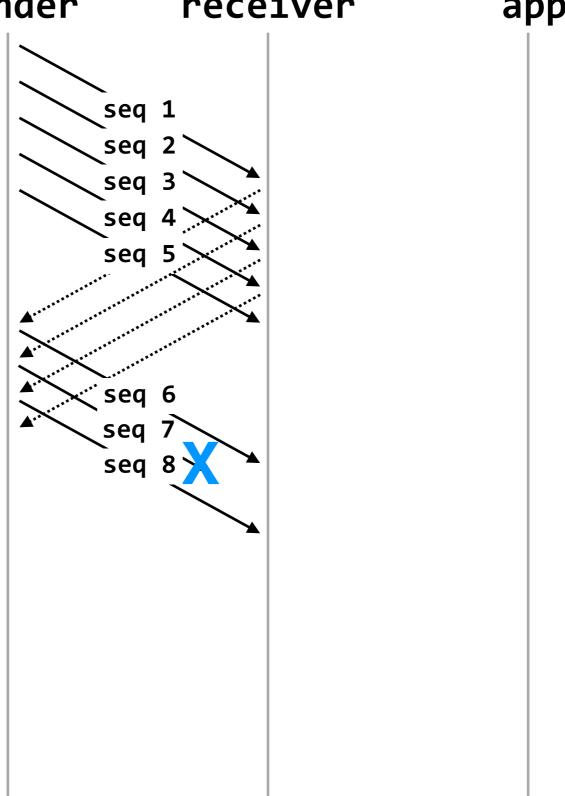
Window: [4,5,6,7,8]



(on same machine)

sender receiver app

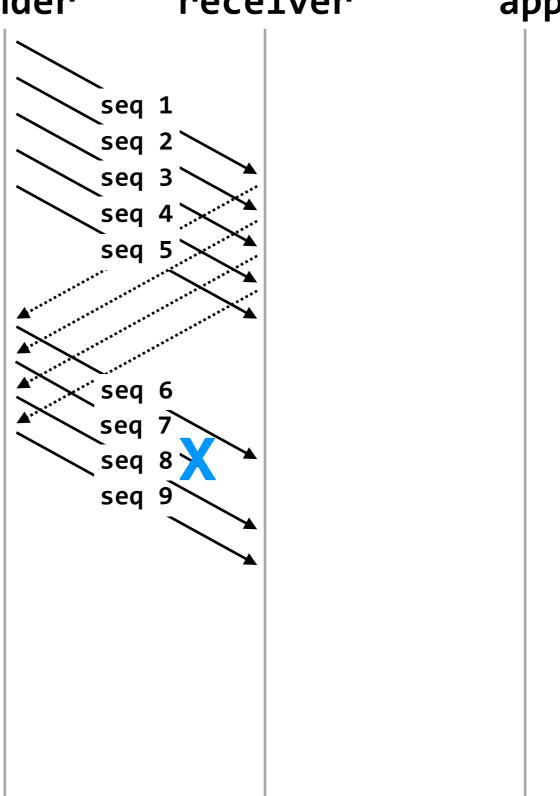
Window: [4,5,6,7,8]



(on same machine)

sender receiver app

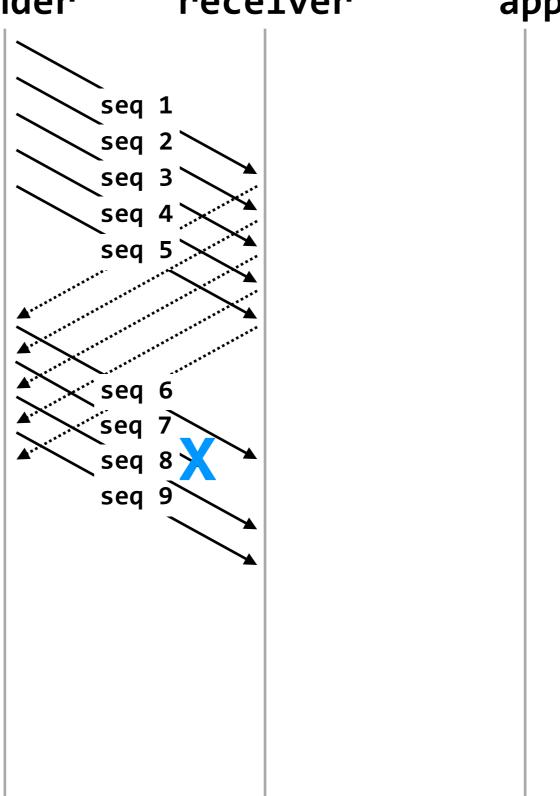
Window: [5,6,7,8,9]



(on same machine)

sender receiver app

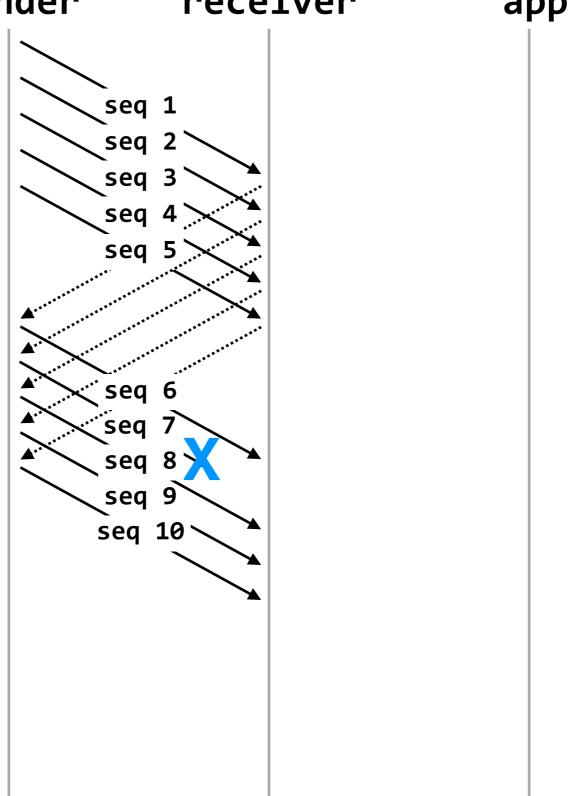
Window: [5,6,7,8,9]



(on same machine)

sender receiver app

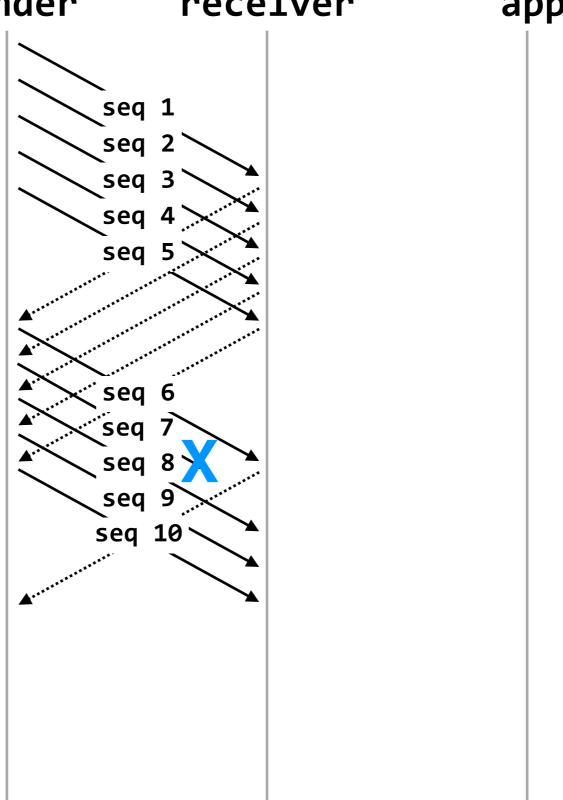
Window: [6,7,8,9,10]



(on same machine)

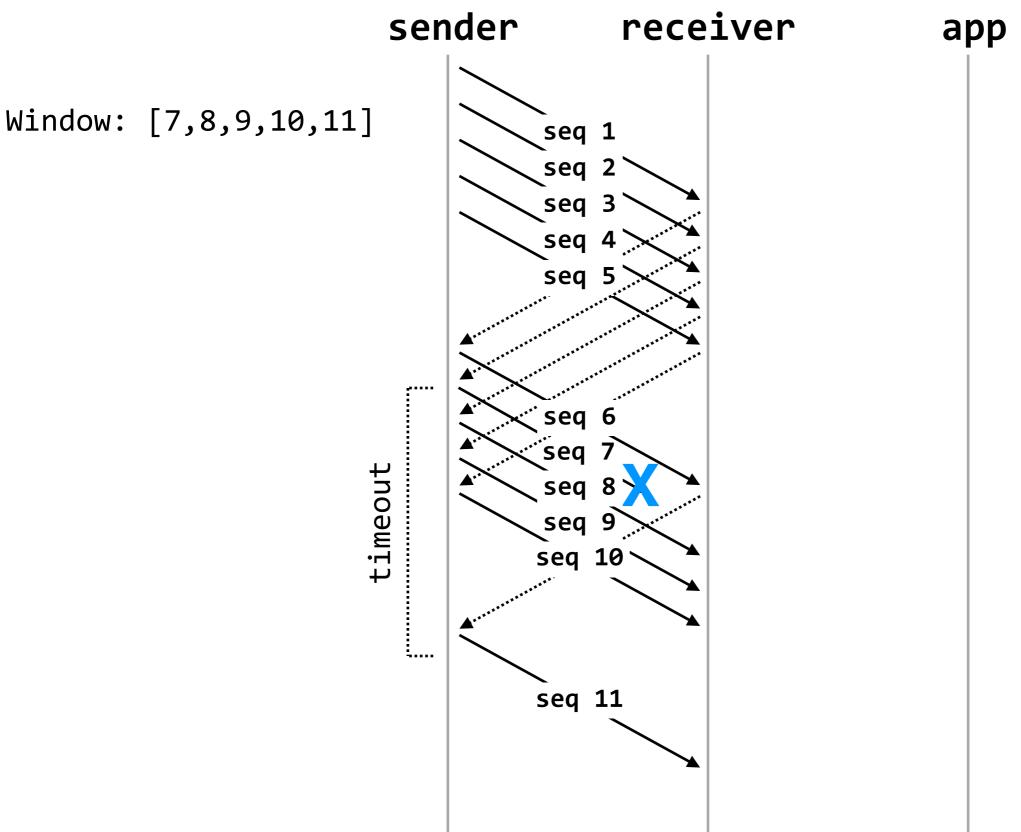
sender receiver app

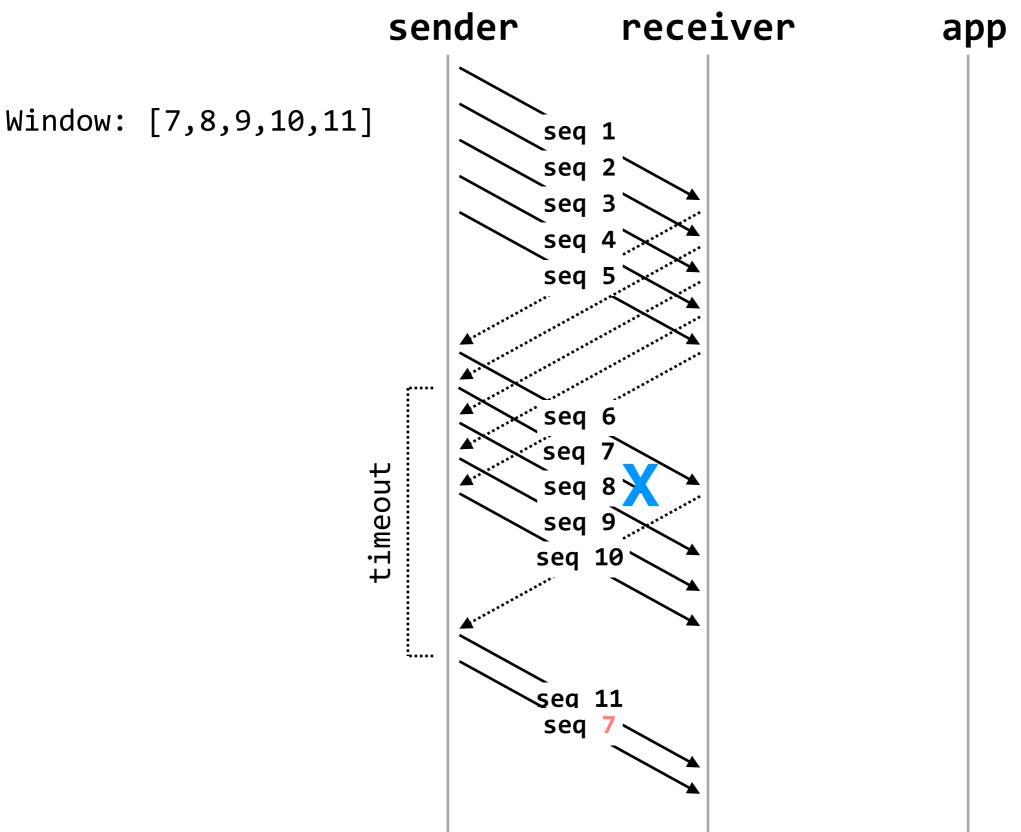
Window: [6,7,8,9,10]

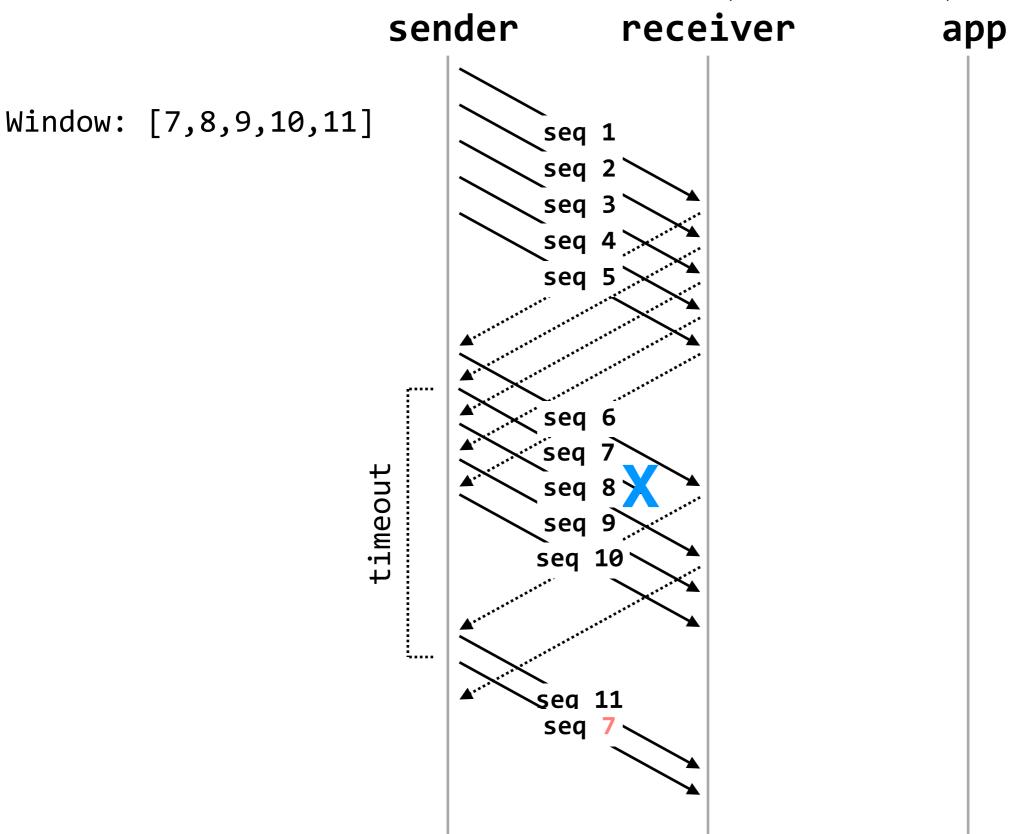


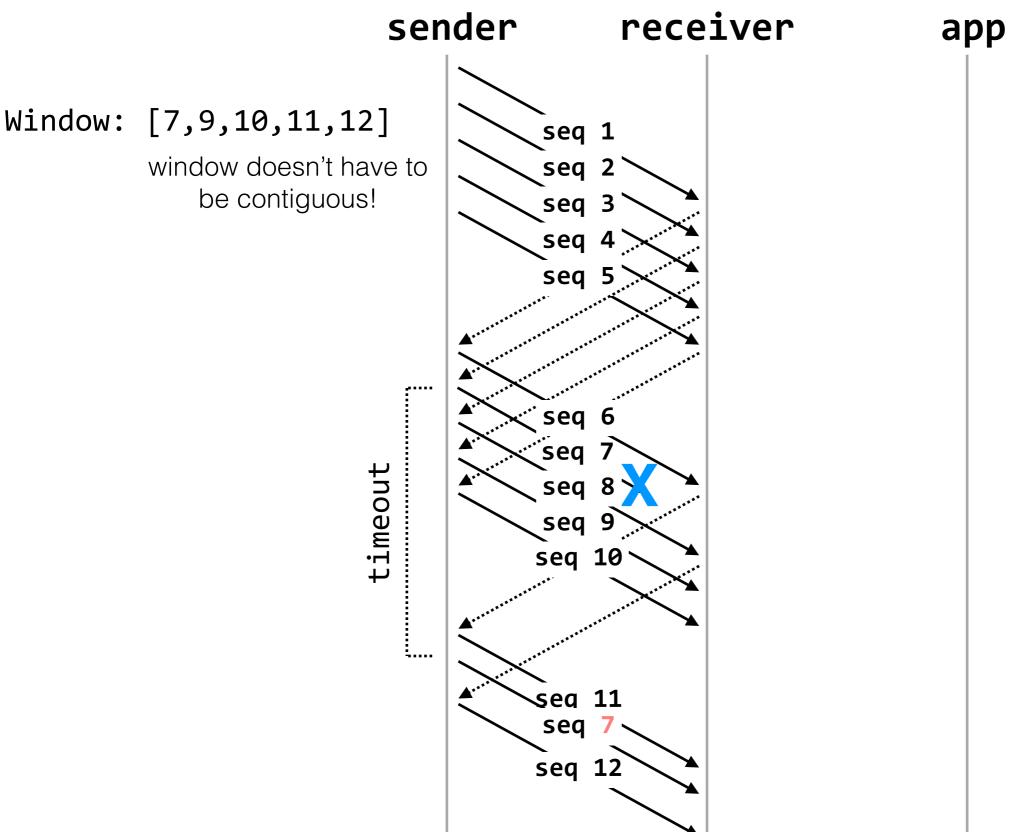
(on same machine) sender receiver app Window: [7,8,9,10,11] seq 1 seq 2 seq seq seq seq seq 8 seq 9 seq 10

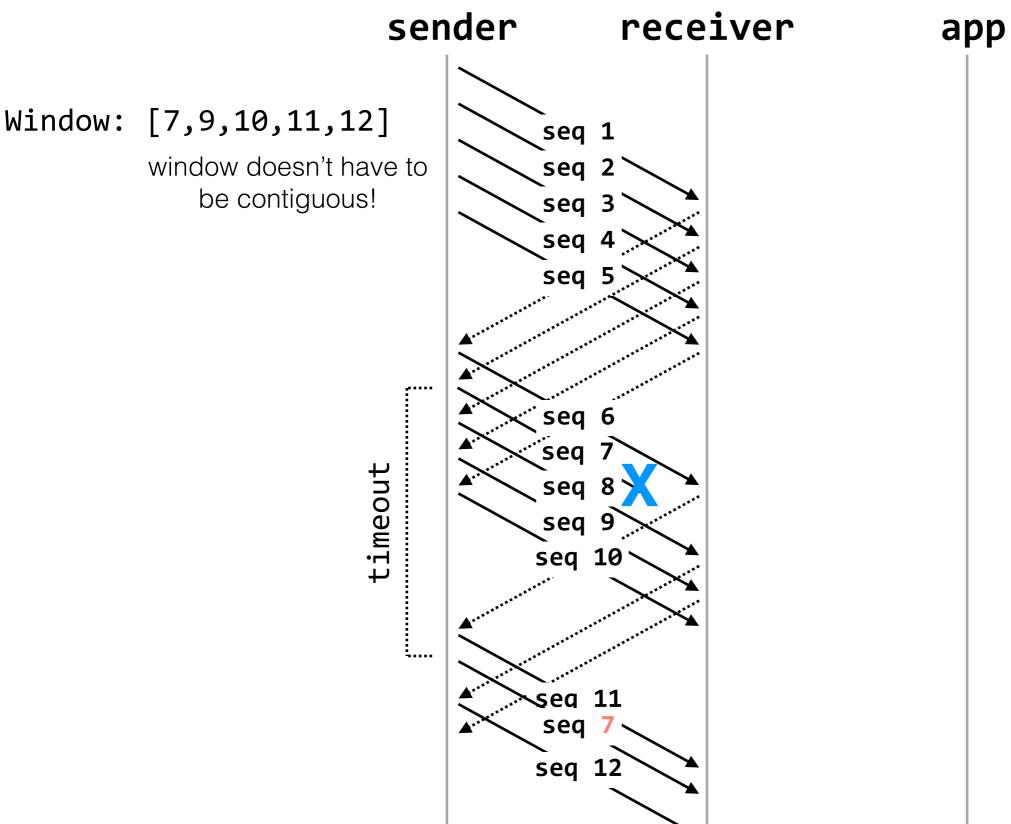
seq 11

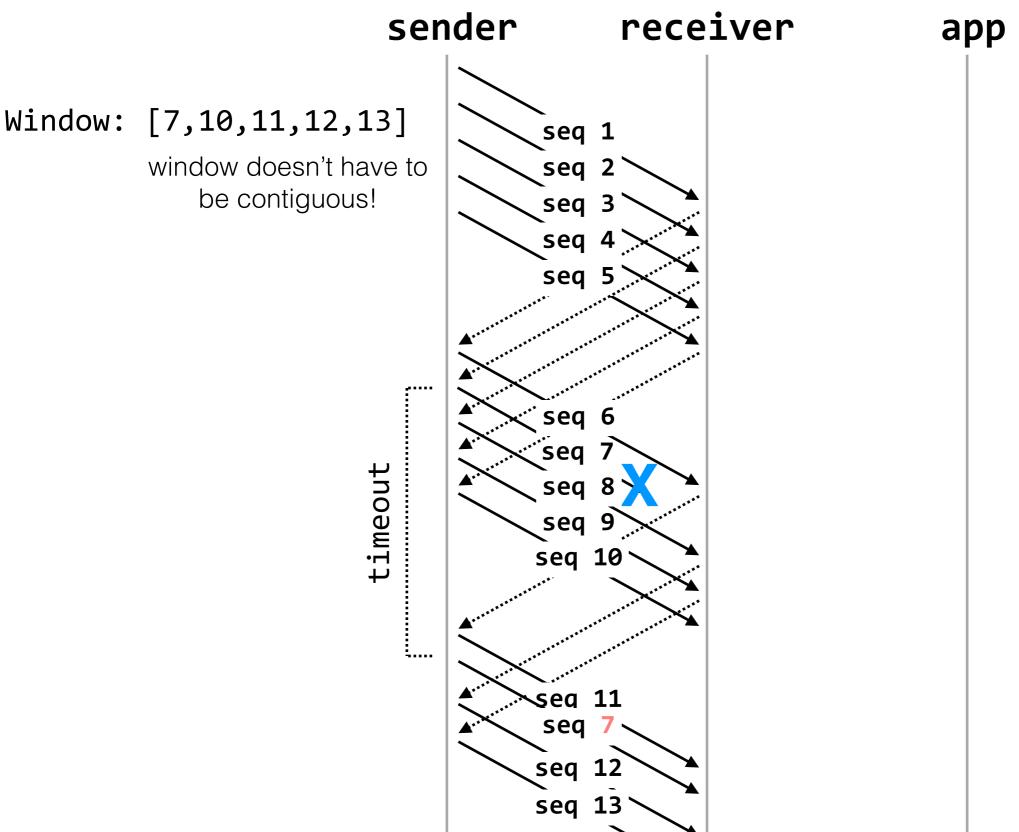


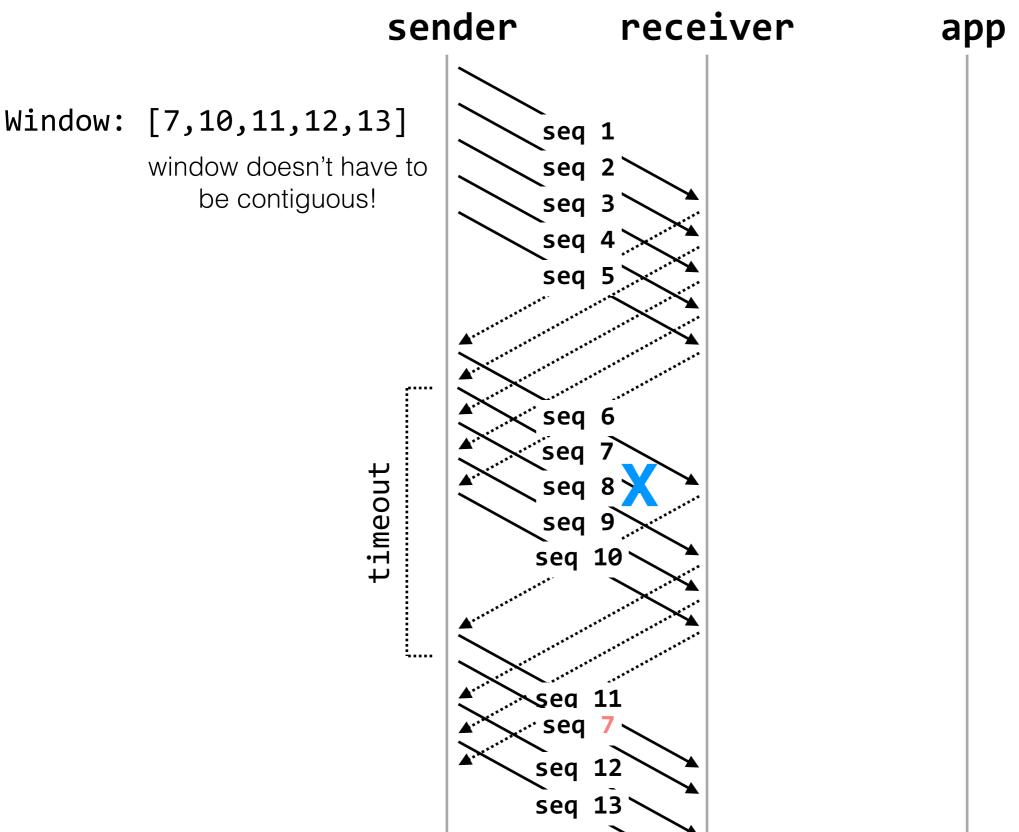


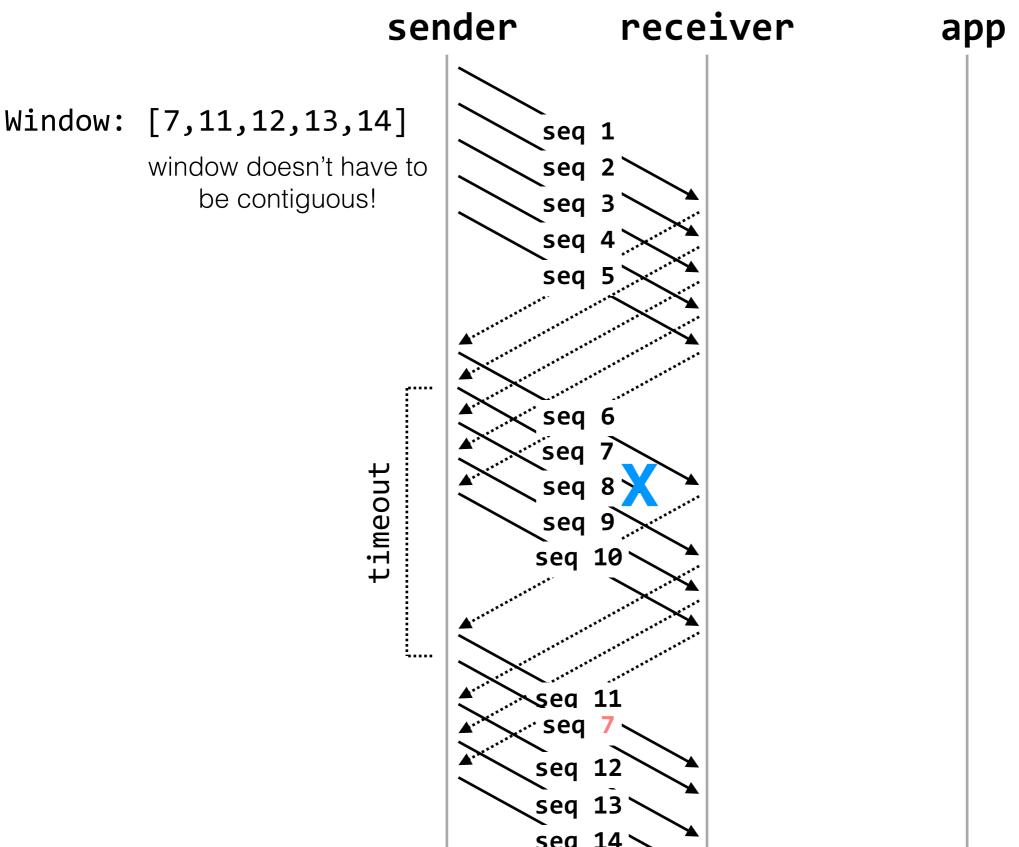


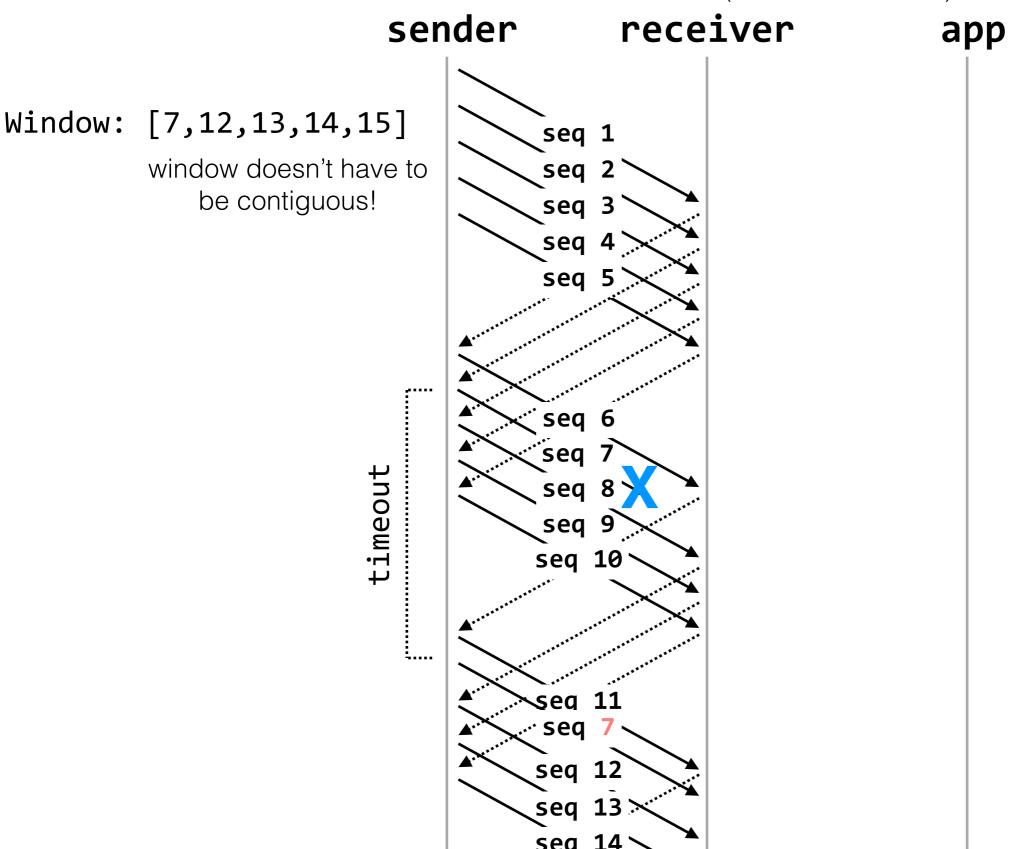


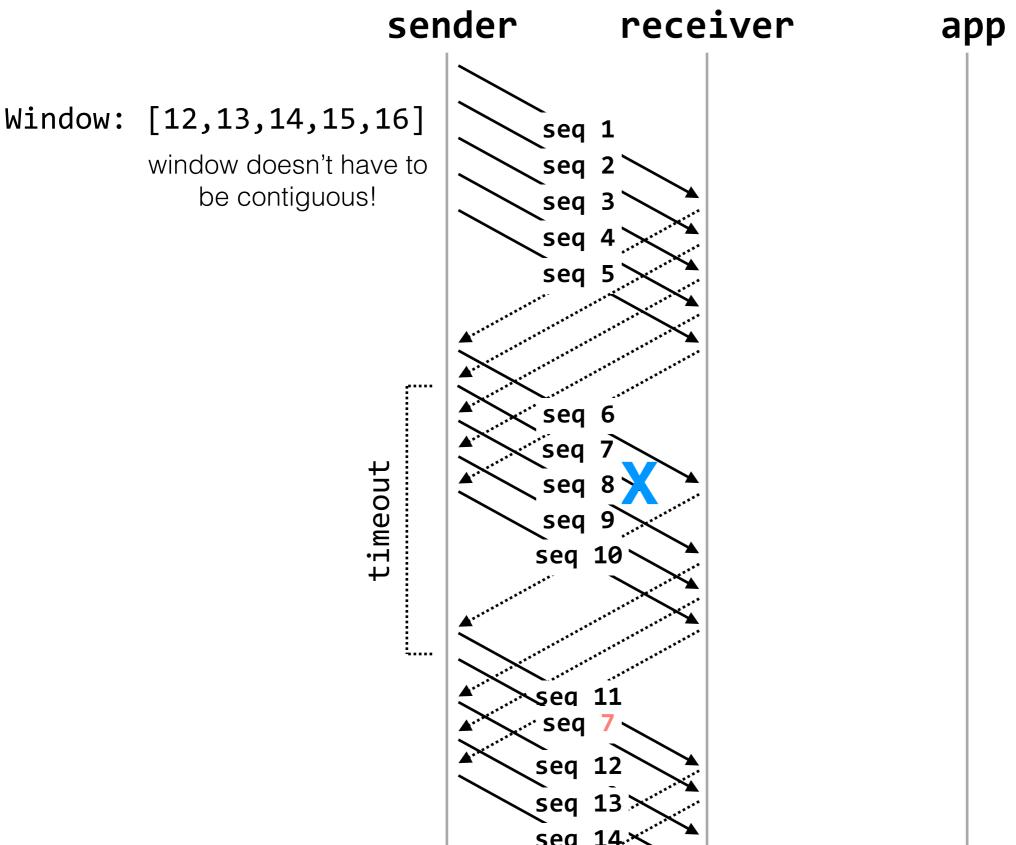










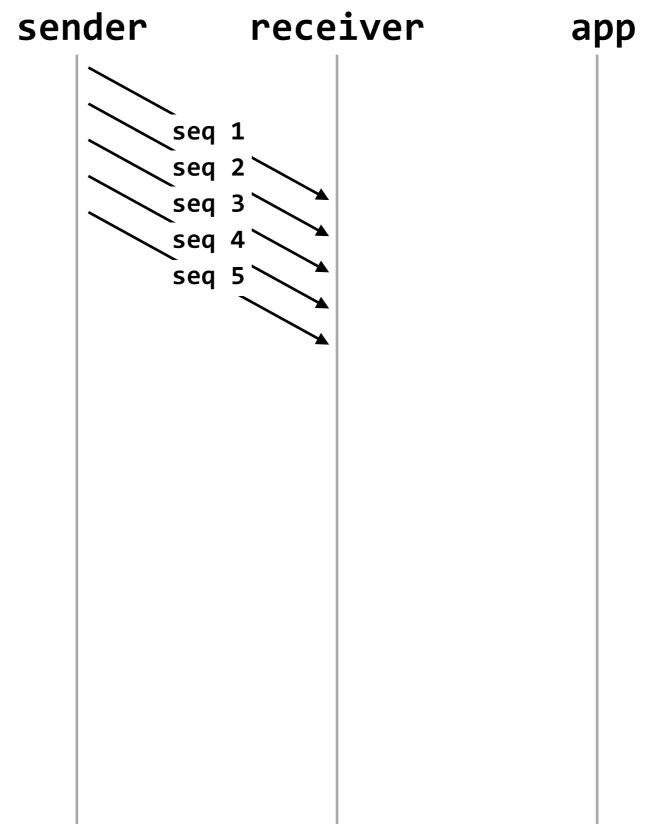


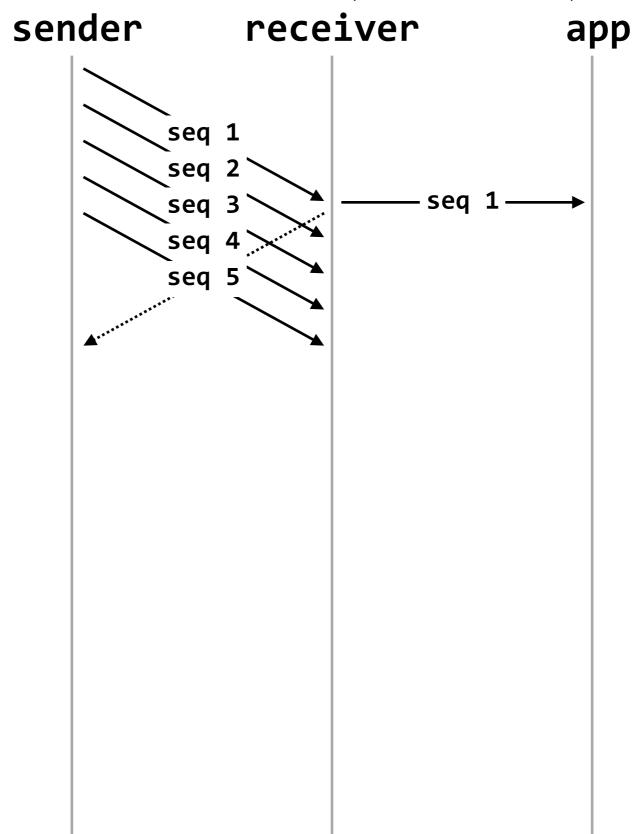
- Transmit a packet if len(un-ACKed list) < W
- Upon transmission of packet k, keep track of k in the un-ACKed list, and the time that k was sent
- When an ACK for packet k is received, remove k from the un-ACKed list.
- Periodically check the un-ACKed list to see if any packets were sent more than timeout seconds ago. If so, re-transmit.

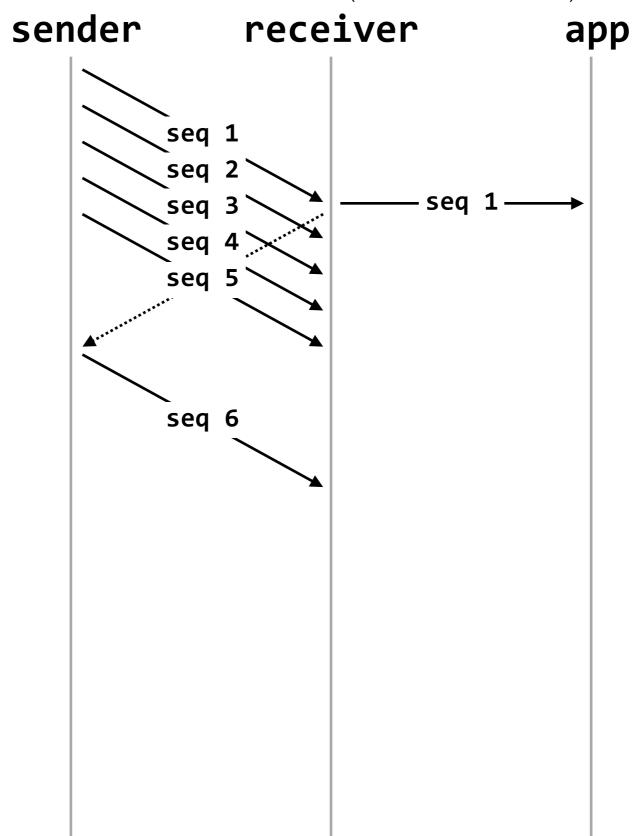
## Sliding-window Protocol: Receiver

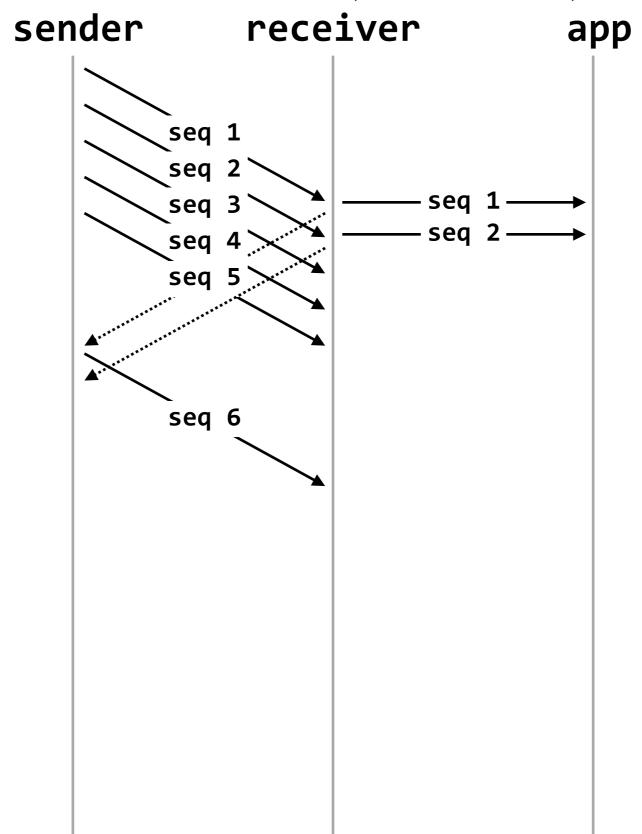
	(OH Same machine)	
sender	receiver	арр

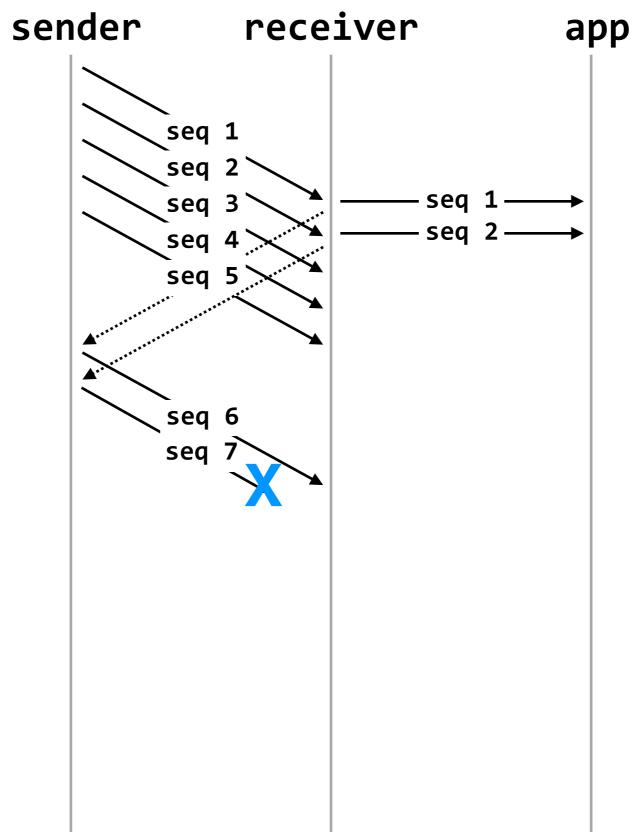
## Sliding-window Protocol: Receiver

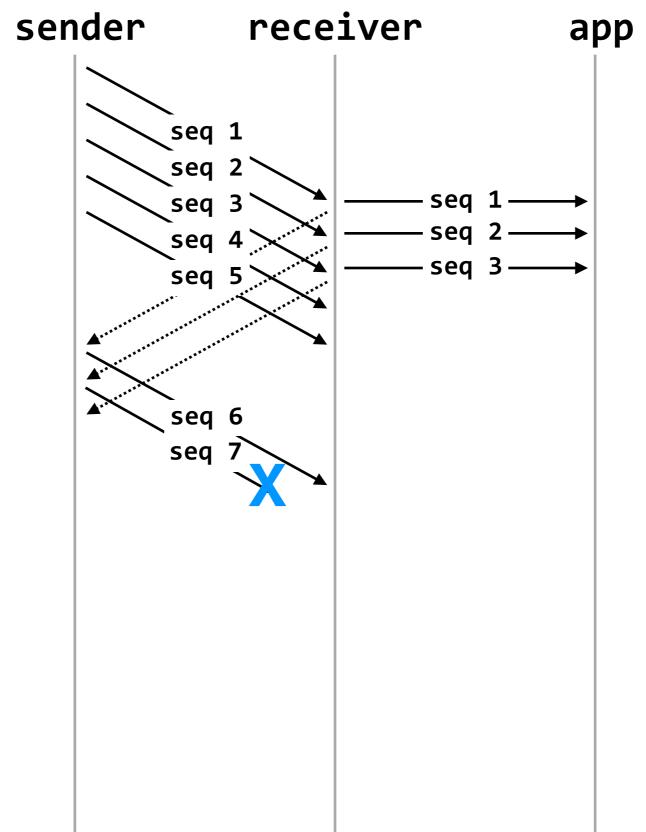


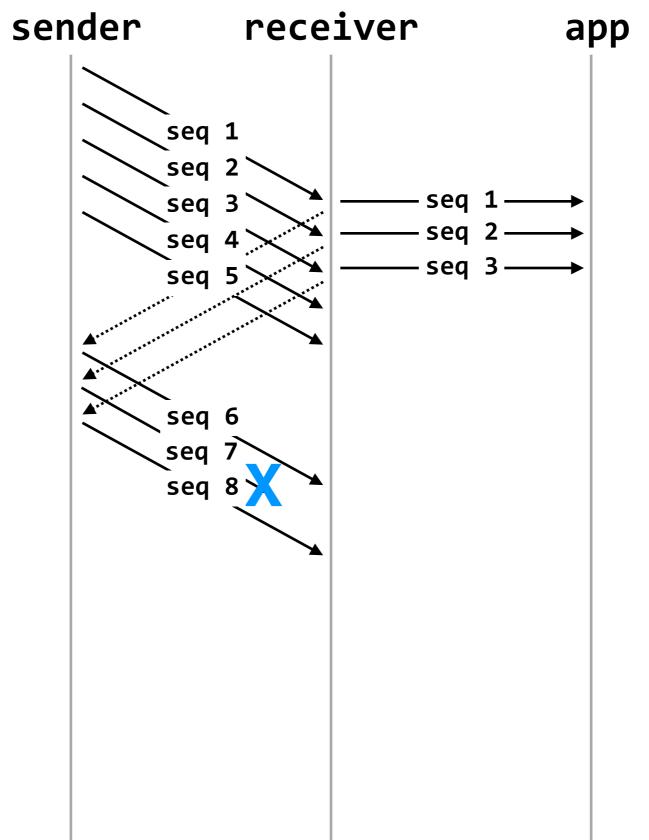


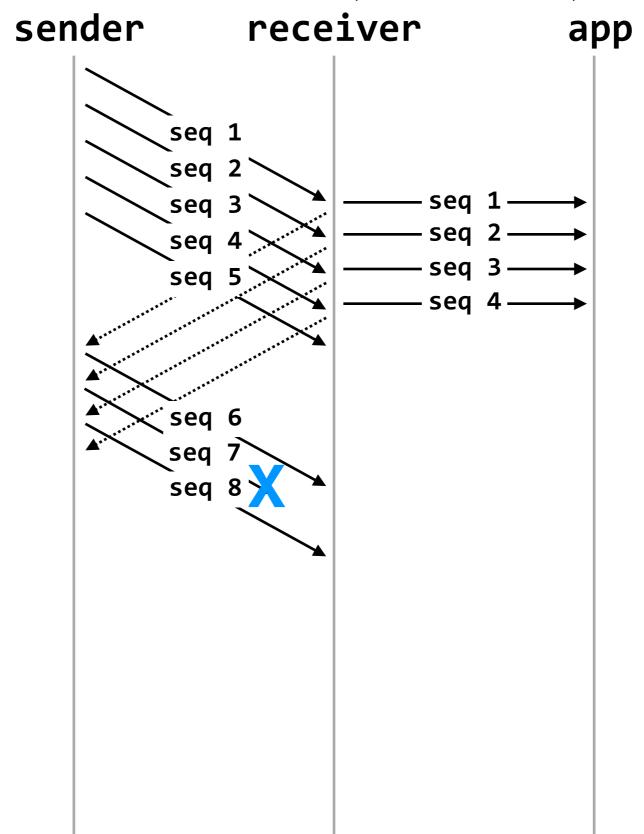


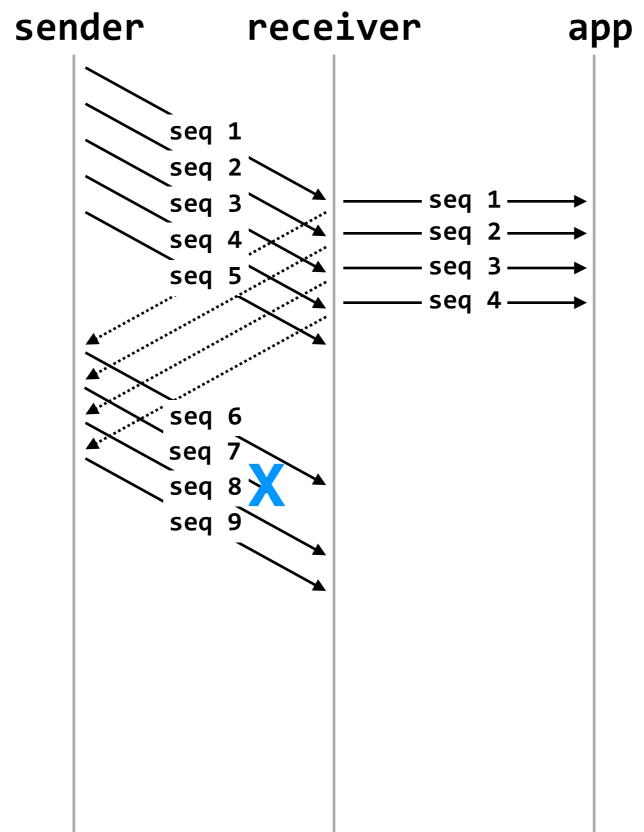


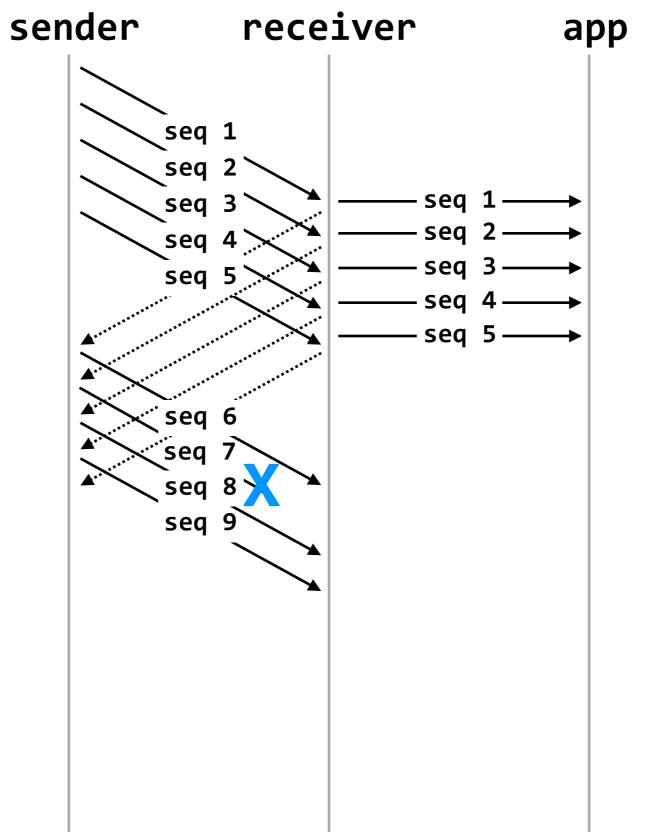


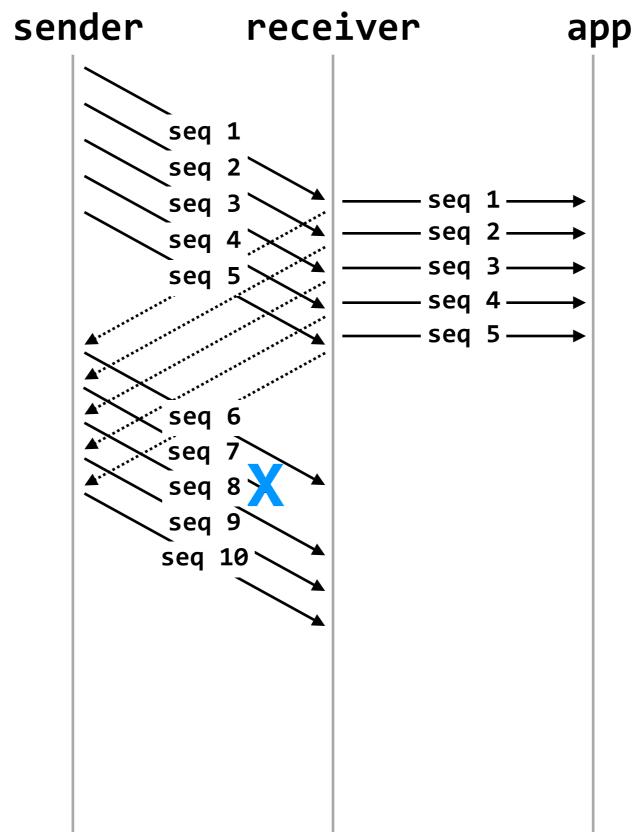


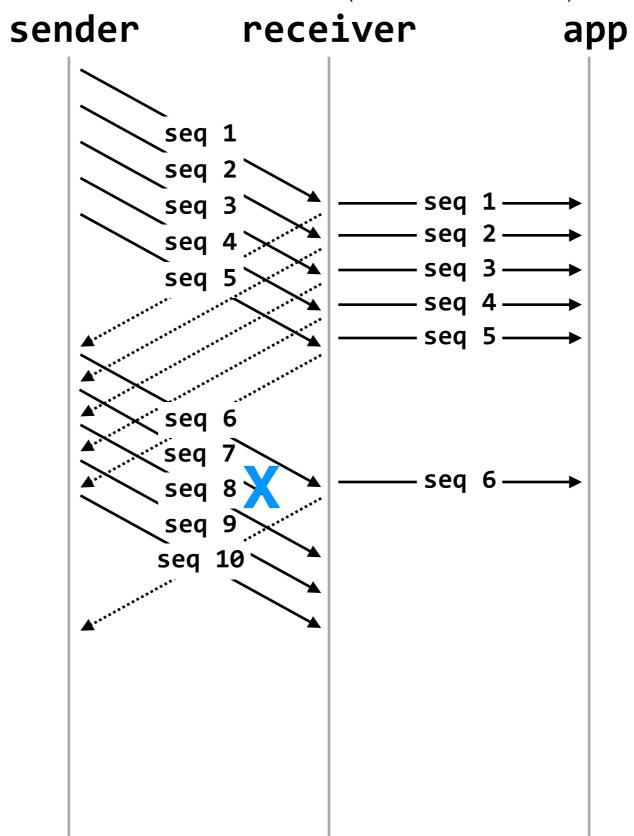


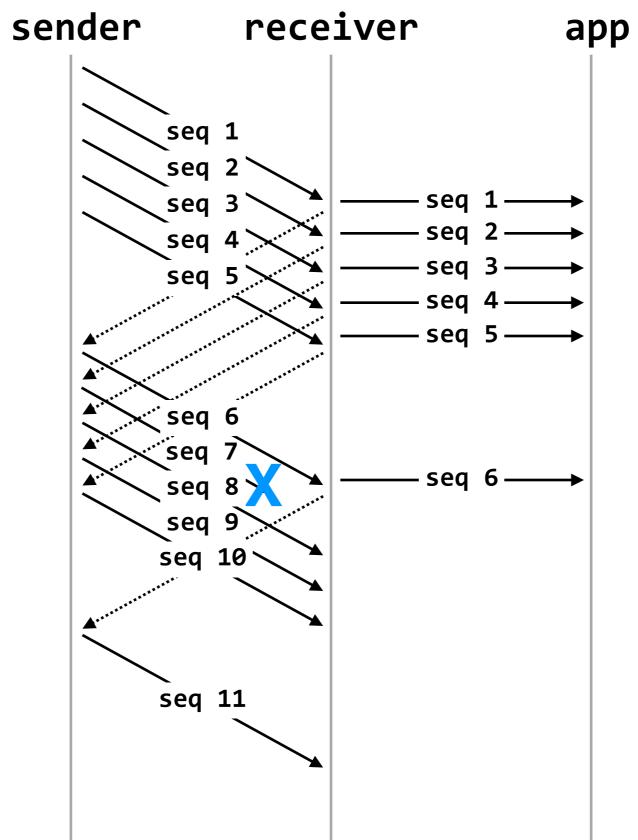


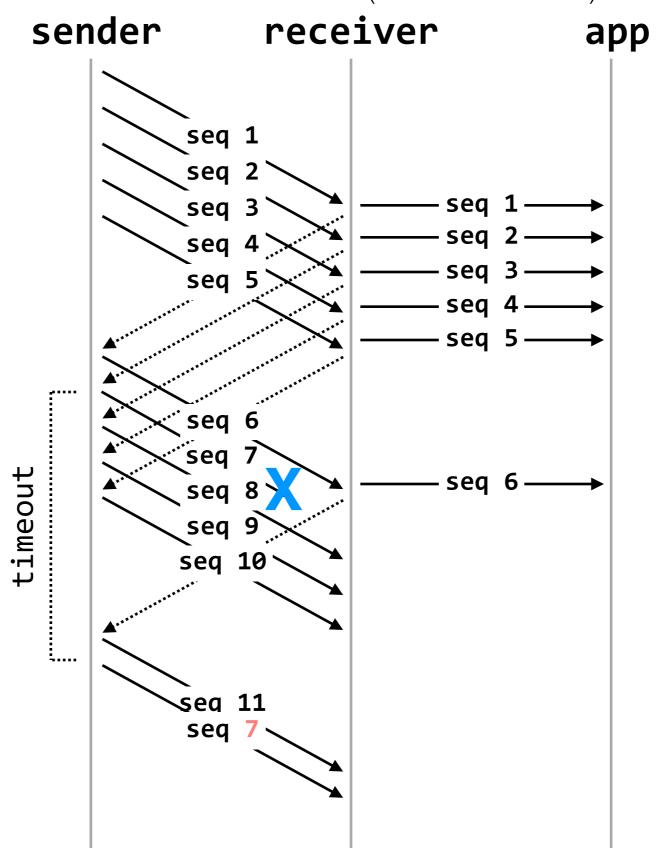


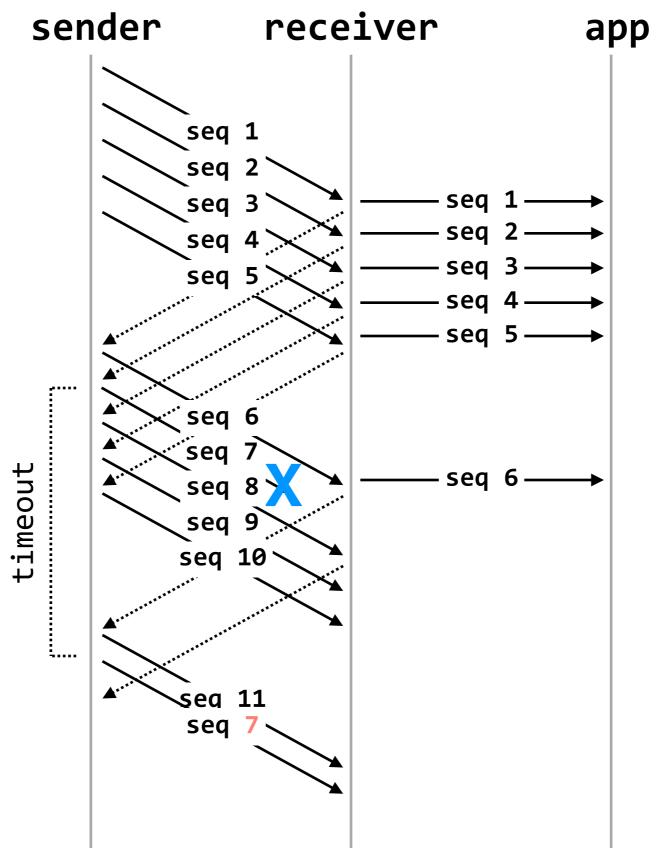


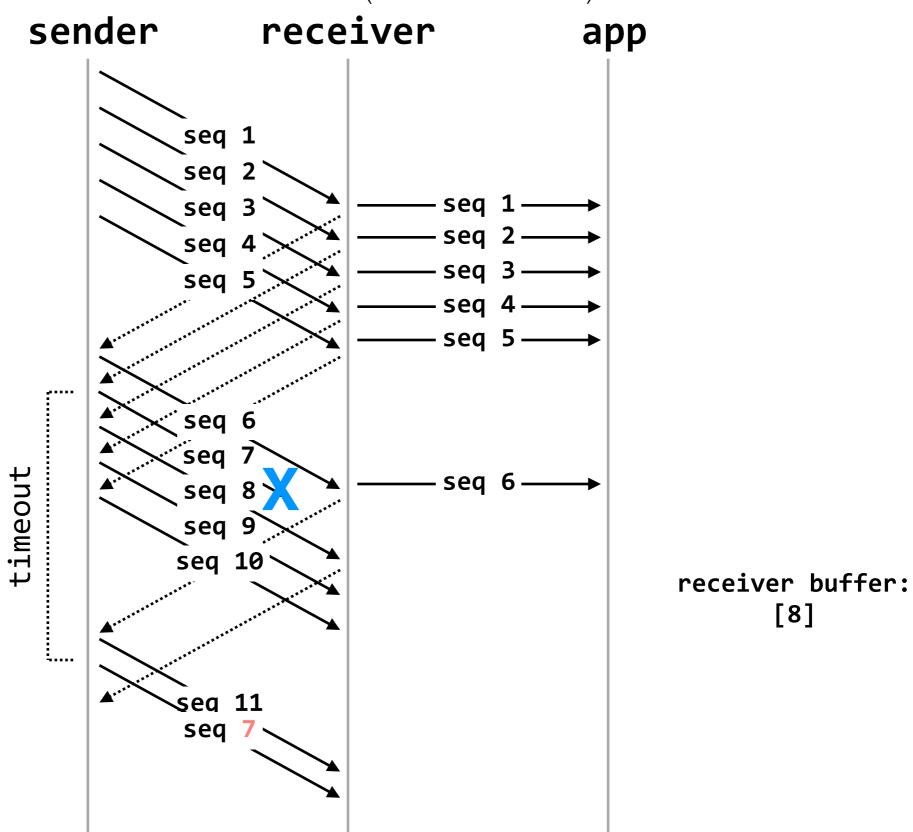


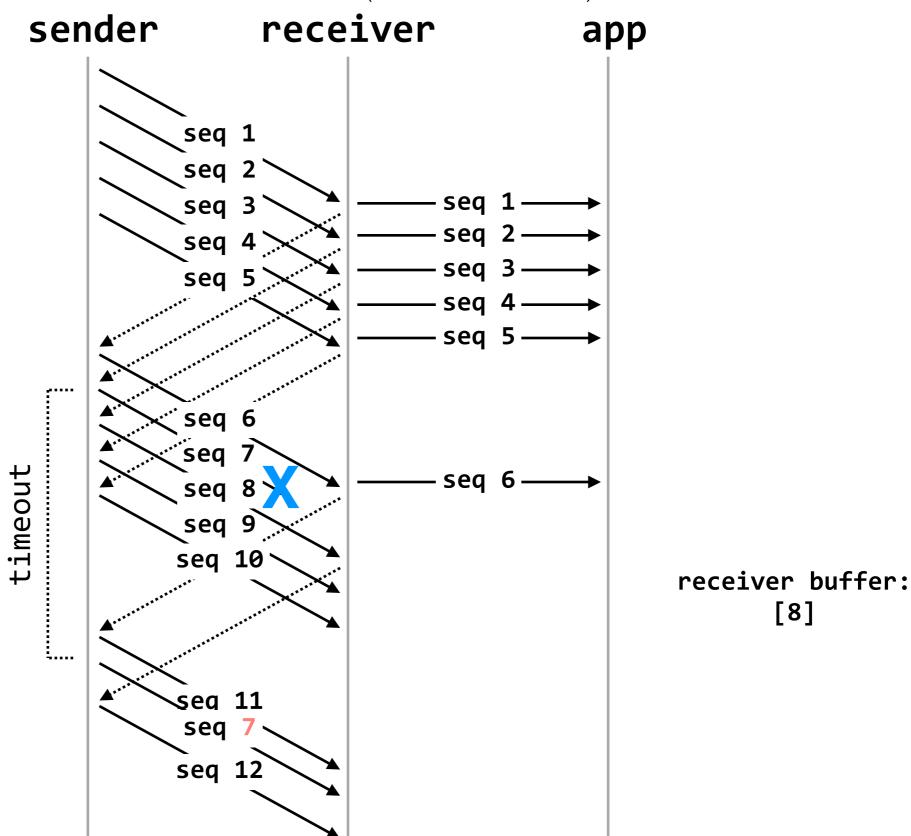


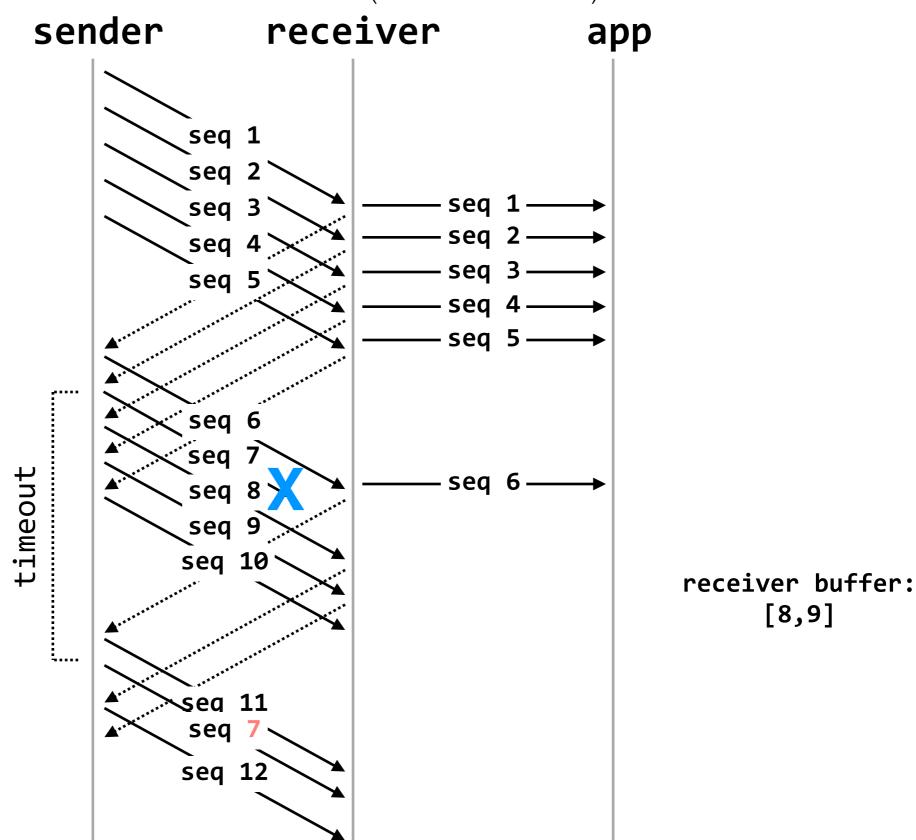


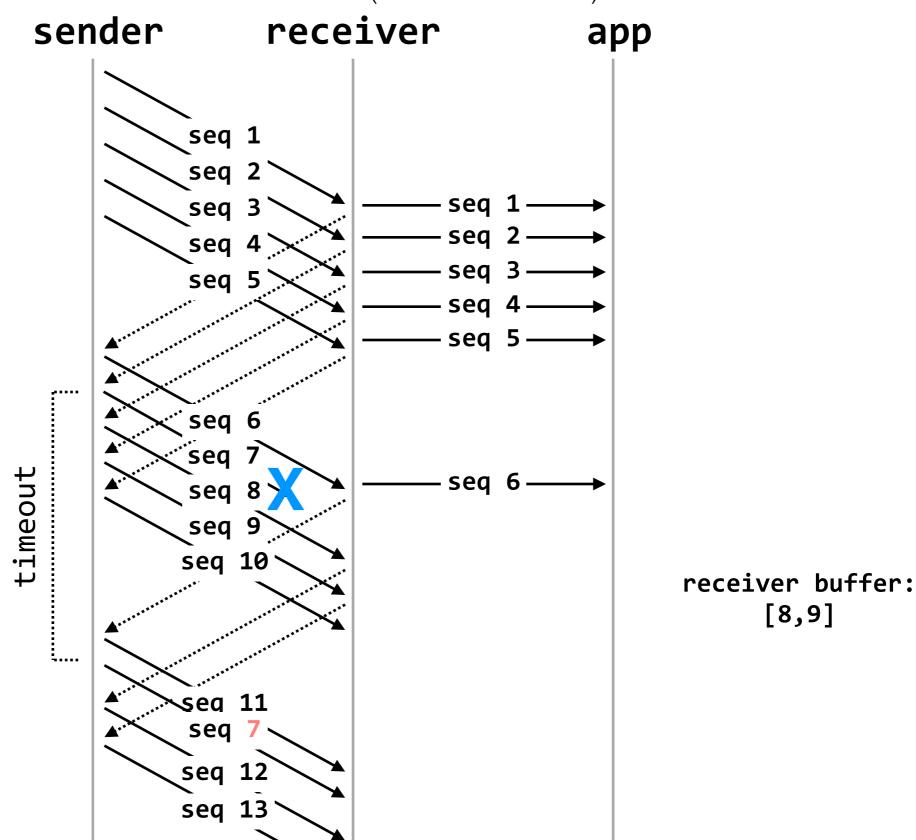


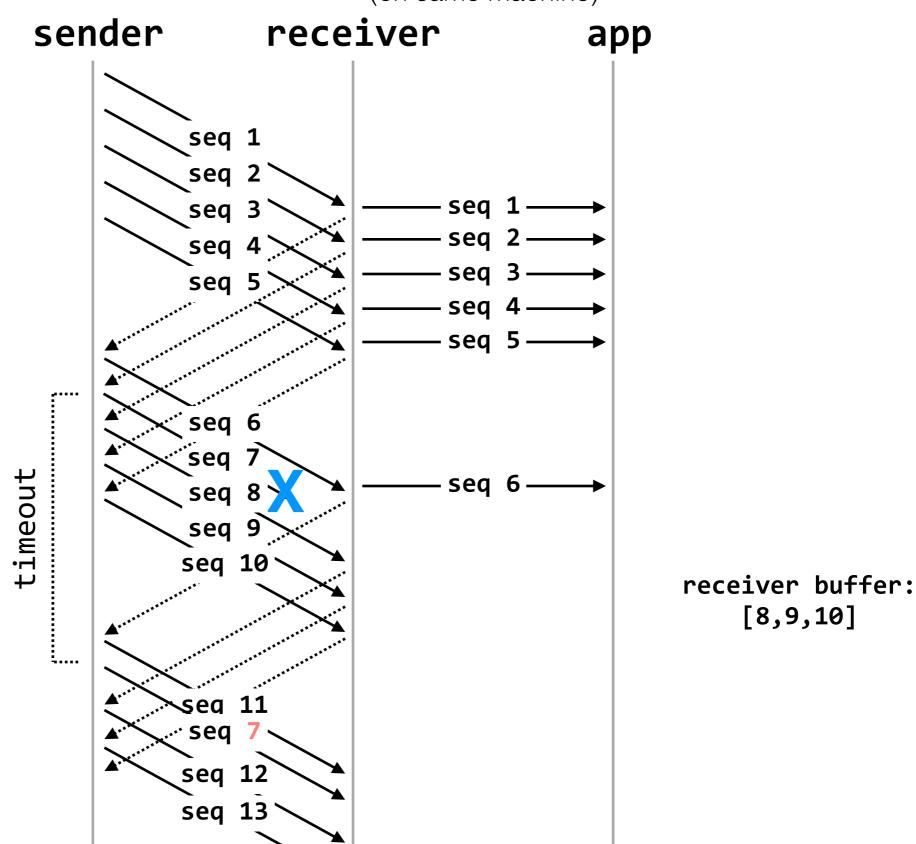


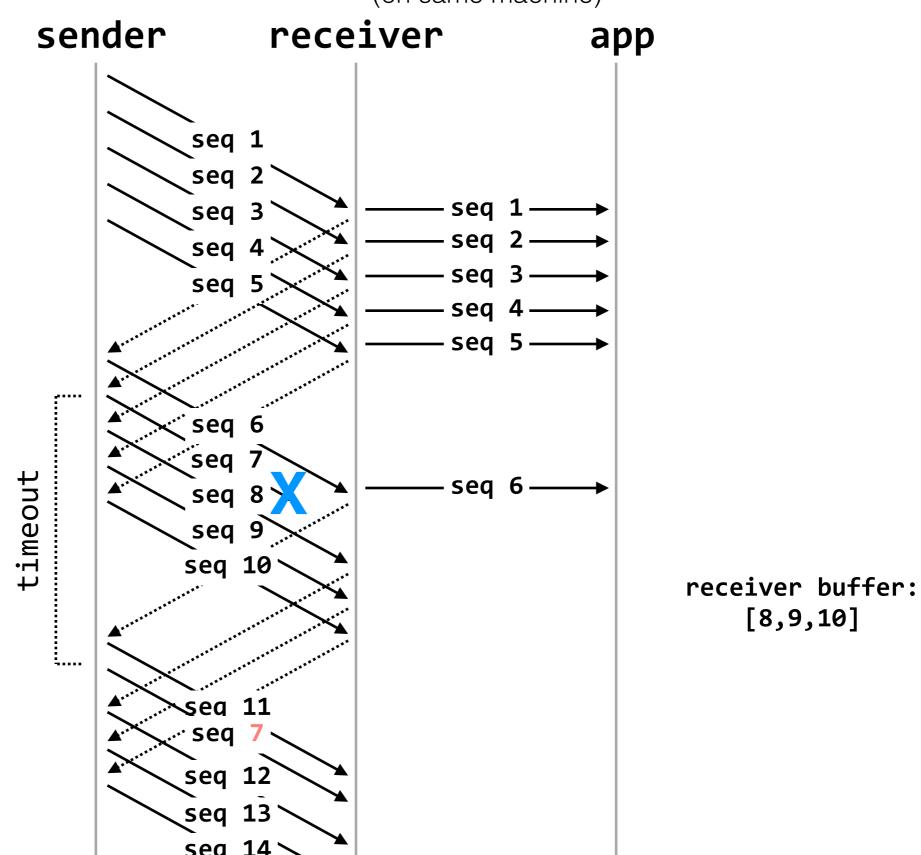


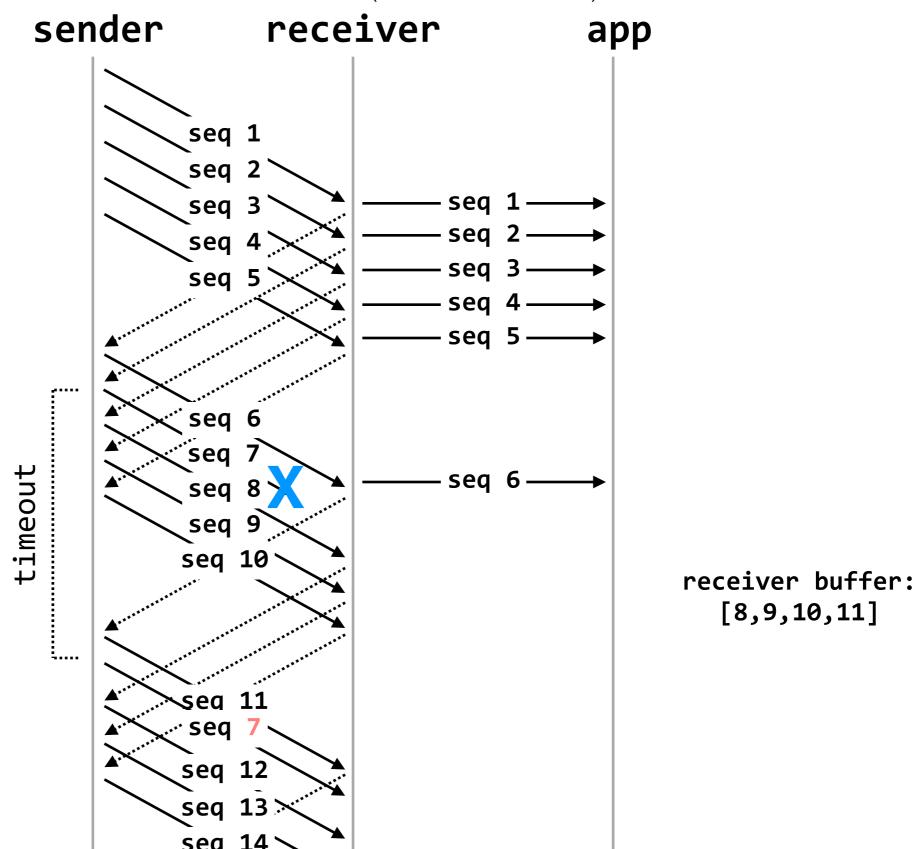


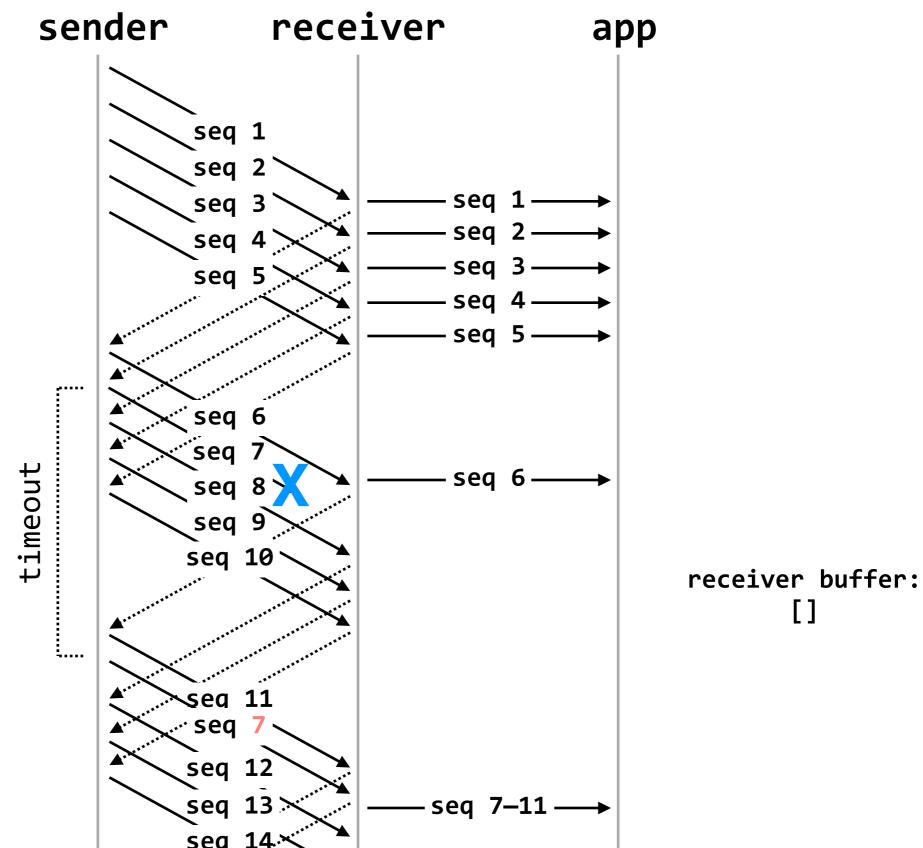




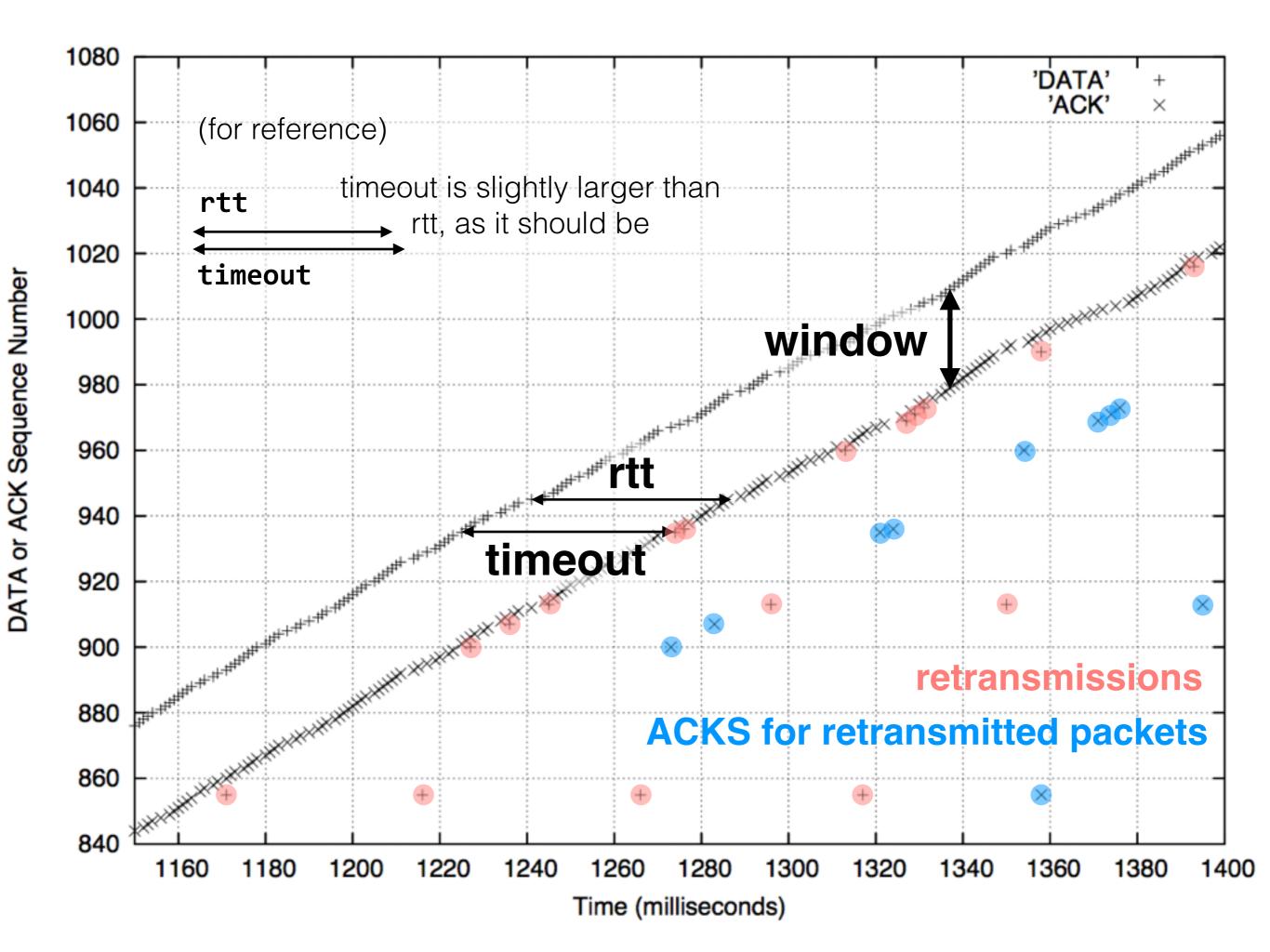




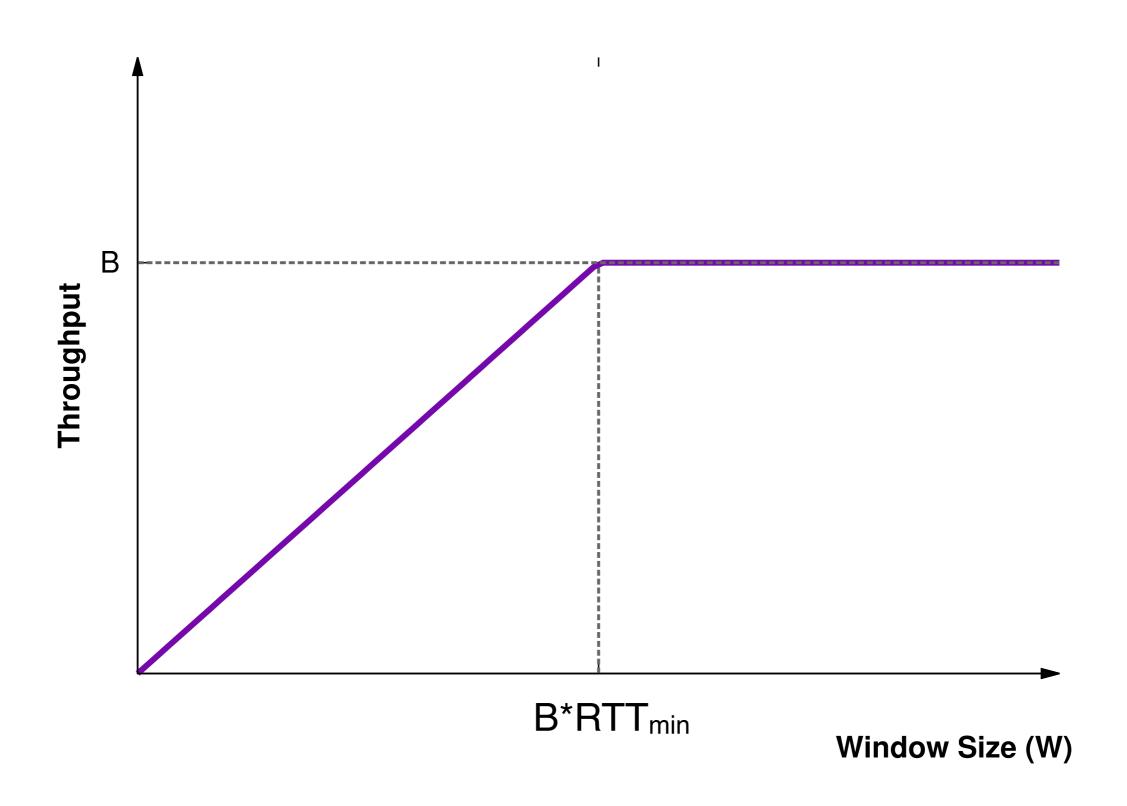


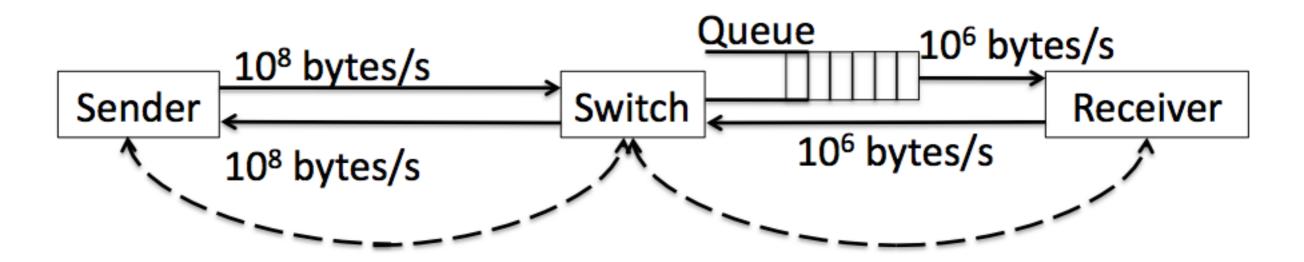


- Send an ACK for every received packet
- Save delivered packets ignoring duplicates in a local buffer
- Keep track of the next packet the application expects. After each reception, deliver as many in-order packets as possible.



## Window Size vs. Throughput





Propagation delay = 0 milliseconds

One-way propagation delay = 10 milliseconds

Max queue size: 100 packets

Packet size: 1000 bytes

ACK size: 40 bytes

Initial window size: 10 packets

- 1. Double W
- 2. Halve the propagation times
- 3. Double bottleneck link rate

# Netflix takes up 9.5% of *upstream* traffic on the North American Internet

ACK packets make Netflix an upload monster during peak viewing hours.

by Jon Brodkin - Nov 20 2014, 7:00am EST





#### Sliding-window protocol

Uses sequence numbers, acknowledgements, and timeouts to ensure exactly-once delivery; allows W packets on the wire at once to improve utilization

#### Setting the window size

W should be at or slightly above (depending on loss) the bandwidth-delay product of the network; this keeps the network utilized without building excessive queues